

# Seminar on Models of Intuitionism

Hand-in exercise 11: model solution

11 May

## Exercise 1.

- (a) Let  $X$  and  $Y$   $G$ -sets. The  $G$ -actions on  $X$  and  $Y$  induce a  $G$ -action on the cartesian product of  $X$  and  $Y$ . Namely, for  $(x, y) \in X \times Y$ , put  $g \cdot (x, y) = (g \cdot x, g \cdot y)$ . One readily verifies that this indeed satisfies the axioms for a group action. So  $X \times Y$  is a  $G$ -set.

We show that the projection  $\pi: X \times Y \rightarrow X$  is a  $G$ -map. Let  $g \in G$  and  $(x, y) \in X \times Y$ . Then  $g \cdot \pi(x, y) = g \cdot x = \pi(g \cdot x, g \cdot y) = \pi(g \cdot (x, y))$ . So  $\pi$  is indeed a  $G$ -map. Similarly, the projection  $\pi': X \times Y \rightarrow Y$  is also a  $G$ -map.

Now suppose we have a  $G$ -set  $Z$  and  $G$ -maps  $a: Z \rightarrow X$  and  $b: Z \rightarrow Y$ . We claim that  $\langle a, b \rangle: Z \rightarrow X \times Y$  given by  $z \mapsto (a(z), b(z))$  is a  $G$ -map. To this end, let  $g \in G$  and  $z \in Z$ . Note that  $g \cdot \langle a, b \rangle(z) = g \cdot (a(z), b(z)) = (g \cdot a(z), g \cdot b(z)) = (a(g \cdot z), b(g \cdot z)) = \langle a, b \rangle(g \cdot z)$ , since  $a$  and  $b$  are  $G$ -maps. Hence,  $\langle a, b \rangle$  is indeed a  $G$ -map.

Also, we immediately see that  $\pi \circ \langle a, b \rangle = a$  and  $\pi' \circ \langle a, b \rangle = b$ . Finally, that  $\langle a, b \rangle$  is the unique map with these properties follows immediately from the corresponding fact in **Set**. ■

*Each paragraph is worth half a point.*

- (b) Suppose  $X$  and  $Y$  are  $G$ -sets. We need to define a  $G$ -action on  $Y^X$  such that the map  $\varepsilon: X \times Y^X \rightarrow Y$  given by  $(x, f) \mapsto f(x)$  is a  $G$ -map. That is:  $g \cdot \varepsilon(x, f) = g \cdot f(x)$  should equal  $\varepsilon(g \cdot (x, f)) = \varepsilon((g \cdot x, g \cdot f)) = (g \cdot f)(g \cdot x)$  for any  $x \in X$  and  $f \in Y^X$ . This suggests putting  $g \cdot f = (x \mapsto g \cdot f(g^{-1} \cdot x))$ . For then, given any  $x \in X$ , we have  $(g \cdot f)(g \cdot x) = g \cdot f(g^{-1} \cdot (g \cdot x)) = g \cdot f((g^{-1}g) \cdot x) = g \cdot f(e \cdot x) = g \cdot f(x)$ , so  $\varepsilon$  is a  $G$ -map.

We verify that this is indeed a group action on  $Y^X$ . Let  $f \in Y^X$  and  $x \in X$ . First of all,  $(e \cdot f)(x) = e \cdot f(e^{-1} \cdot x) = f(x)$ , so  $e \cdot f = f$ . Furthermore, for all  $g, h \in G$ , we have

$$\begin{aligned} (h \cdot (g \cdot f))(x) &= h \cdot ((g \cdot f)(h^{-1} \cdot x)) = h \cdot (g \cdot f(g^{-1} \cdot (h^{-1} \cdot x))) = (hg) \cdot f((g^{-1}h^{-1}) \cdot x) \\ &= (hg) \cdot f((hg)^{-1} \cdot x) = ((hg) \cdot f)(x), \end{aligned}$$

so  $h \cdot (g \cdot f) = (hg) \cdot f$ , as desired.

We proceed by checking the universal property. Let  $Z$  be a  $G$ -set and  $f: X \times Z \rightarrow Y$  a  $G$ -map. We claim that  $\tilde{f}: Z \rightarrow Y^X$  given by  $z \mapsto (x \mapsto f(x, z))$  is a  $G$ -map. Let  $g \in G$  and  $z \in Z$  be arbitrary. Note that since  $f$  is a  $G$ -map, we have for all  $x \in X$  that:

$$\left( g \cdot \tilde{f}(z) \right) (x) = g \cdot f(g^{-1} \cdot x, z) = f(g \cdot (g^{-1} \cdot x), g \cdot z) = f(e \cdot x, g \cdot z) = f(x, g \cdot z) = \tilde{f}(g \cdot z)(x),$$

so  $g \cdot \tilde{f}(z) = \tilde{f}(g \cdot z)$ . Thus,  $\tilde{f}$  is indeed a  $G$ -map.

Finally, we know that  $\varepsilon \circ \langle \pi, \tilde{f} \circ \pi' \rangle = f$  as we have seen this in **Set**. That  $\tilde{f}$  is the unique arrow with this property follows from the corresponding fact in **Set**. ■

*One point each for the first two paragraphs. Half a point for each of the final two paragraphs.*

**Exercise 2.**

(4) Let  $A \stackrel{f}{\Rightarrow} \top$  be a deduction. The equation  $f = !_A$  identifies  $A \stackrel{f}{\Rightarrow} \top$  with  $\overline{A \Rightarrow \top}^{\text{TRUE}}$ . In other words, there is, up to equivalence, exactly one deduction of the entailment  $A \Rightarrow \top$ , namely the one given by the TRUE-rule. ■

(5) Let  $C \stackrel{f}{\Rightarrow} A$  and  $C \stackrel{g}{\Rightarrow} B$  be deductions. The equation  $\pi \circ \langle f, g \rangle = f$  identifies

$$\frac{\frac{C \stackrel{f}{\Rightarrow} A \quad C \stackrel{g}{\Rightarrow} B}{C \Rightarrow A \wedge B} \wedge_R \quad \overline{A \wedge B \Rightarrow A}^{\wedge L1}}{C \Rightarrow A} \text{CUT}$$

with  $C \stackrel{f}{\Rightarrow} A$  itself. ■

(7) Let  $C \stackrel{h}{\Rightarrow} A \wedge B$  be a deduction. The equation  $\langle \pi \circ h, \pi' \circ h \rangle = h$  identifies

$$\frac{\frac{C \stackrel{h}{\Rightarrow} A \wedge B \quad \overline{A \wedge B \Rightarrow A}^{\wedge L1}}{C \Rightarrow A} \text{CUT} \quad \frac{C \stackrel{h}{\Rightarrow} A \wedge B \quad \overline{A \wedge B \Rightarrow B}^{\wedge L2}}{C \Rightarrow B} \text{CUT}}{C \Rightarrow A \wedge B} \wedge_R$$

with  $C \stackrel{h}{\Rightarrow} A \wedge B$  itself. ■

(10) Let  $A \vee B \stackrel{h}{\Rightarrow} C$  be a deduction. The equation  $[h \circ \kappa, h \circ \kappa'] = h$  identifies

$$\frac{\frac{\overline{A \Rightarrow A \vee B}^{\vee R1} \quad A \vee B \stackrel{h}{\Rightarrow} C}{A \Rightarrow C} \text{CUT} \quad \frac{\overline{B \Rightarrow A \vee B}^{\vee R2} \quad A \vee B \stackrel{h}{\Rightarrow} C}{B \Rightarrow C} \text{CUT}}{A \vee B \Rightarrow C} \vee_L$$

with  $A \vee B \stackrel{h}{\Rightarrow} C$  itself. ■

(12) Let  $C \stackrel{k}{\Rightarrow} A \rightarrow B$  be a deduction. The equation  $(\varepsilon \circ \langle \pi, k \circ \pi' \rangle)^\sim = k$  identifies

$$\frac{\frac{\overline{A \wedge C \Rightarrow A}^{\wedge L1} \quad \frac{\overline{A \wedge C \Rightarrow C}^{\wedge L2} \quad C \stackrel{k}{\Rightarrow} A \rightarrow B}{A \wedge C \Rightarrow A \rightarrow B} \wedge_R}{A \wedge C \Rightarrow A \wedge (A \rightarrow B)} \text{CUT} \quad \frac{\overline{A \wedge (A \rightarrow B) \Rightarrow B}^{\rightarrow L}}{A \wedge (A \rightarrow B) \Rightarrow B} \text{CUT}}{\frac{A \wedge C \Rightarrow B}{C \Rightarrow A \rightarrow B} \rightarrow_R} \text{CUT}$$

with  $C \stackrel{k}{\Rightarrow} A \rightarrow B$  itself. ■