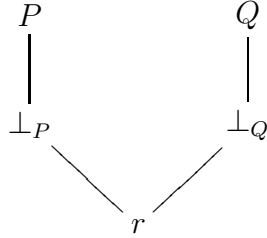


Exercises

Proof Theory, 06-04-2011

Exercise Let P and Q be partial orders with least elements \perp_P and \perp_Q respectively. We define the partial order $P \oplus Q$ as the disjoint union of P and Q with a new least element r added:



Suppose $\{D(p) \mid p \in P\}$, $\{D(q) \mid q \in Q\}$, $\{f_{pp'} \mid p \leq p' \text{ in } P\}$, $\{f_{qq'} \mid q \leq q' \text{ in } Q\}$, $\{[g]_p \mid g \in \mathcal{L}, p \in P\}$, etcetera, make P and Q into Kripke structures for a language \mathcal{L} .

Define: $D(r) = D(\perp_P) \times D(\perp_Q)$, $f_{r\perp_P}$ and $f_{r\perp_Q}$ are the first and second projections. For an n -ary function symbol g let

$$[g]_r((\xi_1, \eta_1), \dots, (\xi_n, \eta_n)) = ([g]_{\perp_P}(\xi_1, \dots, \xi_n), [g]_{\perp_Q}(\eta_1, \dots, \eta_n))$$

and for an n -ary relation symbol R define: $((\xi_1, \eta_1), \dots, (\xi_n, \eta_n)) \in [R]_r$ if and only if both $(\xi_1, \dots, \xi_n) \in [R]_{\perp_P}$ and $(\eta_1, \dots, \eta_n) \in [R]_{\perp_Q}$.

- a. Show that this definition makes $P \oplus Q$ into a Kripke structure for \mathcal{L} .
- b. Show that for every atomic formula $\phi(a_1, \dots, a_n)$ and n -tuple $((\xi_1, \eta_1), \dots, (\xi_n, \eta_n))$ of elements of $D(r)$, we have: $r \Vdash \phi[(\xi_1, \eta_1), \dots, (\xi_n, \eta_n)]$ if and only if both $\perp_P \Vdash \phi[(\xi_1, \dots, \xi_n)]$ and $\perp_Q \Vdash \phi[(\eta_1, \dots, \eta_n)]$.
- c. Show that the equivalence of part b extends to all Harrop formulas.
- d. Use c to give a Kripke-model theoretic proof of the statement that if the sequent $\Gamma \rightarrow A \vee B$ is LJ-provable, where Γ consists of Harrop formulas, then at least one of the sequents $\Gamma \rightarrow A$, $\Gamma \rightarrow B$ is LJ-provable.

[Hint: use the completeness theorem for Kripke models]